

# Presentation Slides

## Chapter 9

### Multi-value Numeric Design

#### Logically Determined Design:

Clockless System Design With NULL Convention Logic

by Karl Fant

John Wiley & Sons, Inc.

Introduce quaternary logic

# Does Multi-value Representation Provide Advantages For Expressing Numeric Processes?

With the familiar electronic representation a digit is a path and value is represented as discriminable voltage levels on that path; values are clearly much more expensive than digits.

With multi-path representation enabling multi-value variables, values and digits are both represented with paths. The cost to represent a digit and the cost to represent its values are identical and there is no bias.

Multi-path representation allows the consideration of radices and encodings other than binary for numeric representation.

Various encodings are considered in terms of cost of transmission and cost of functional combination.

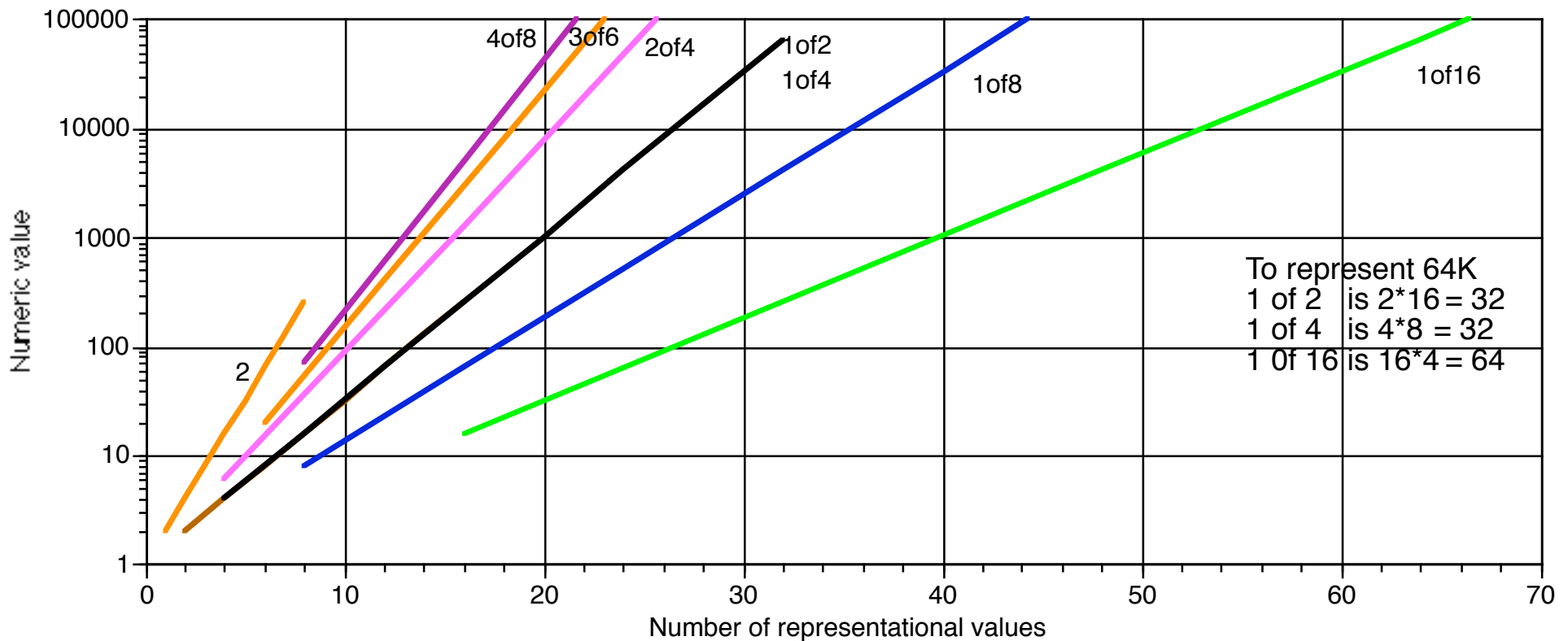
# Cost of Representation

Cost of place  
Cost of value

Binary  
voltage threshold  
wire

Multi-rail  
wire  
wire

The cost of representing a given number with M of N encoding is the number of values per digit times the number of digits to represent the number.



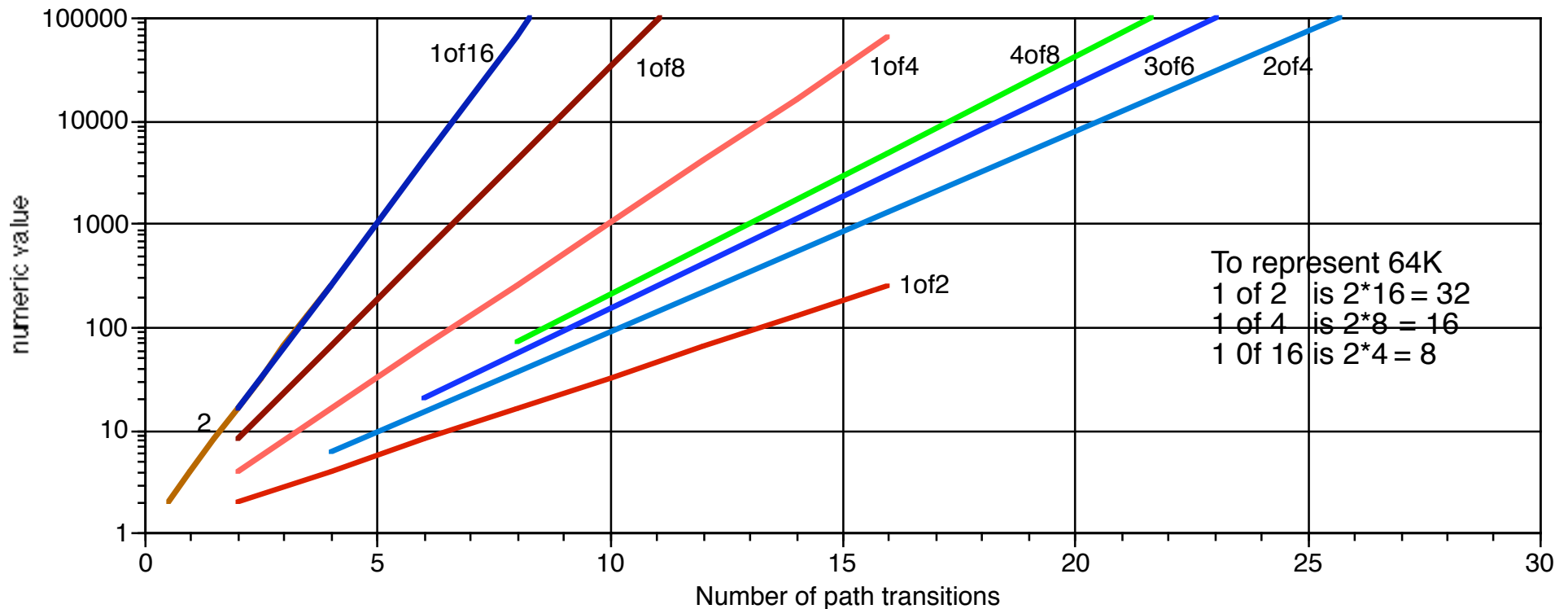
The N/2 of N encodings show significant advantage.

The 1 of 2 and 1 of 4 encodings are the optimal 1 of N encodings.

# Cost of Flow

	binary	$N/2 \text{ of } N$	$1 \text{ of } N$
Cost of flow per digit	0.5	$N$	2

The cost of flow is how many path switches per digit times the number of digits to move a number.



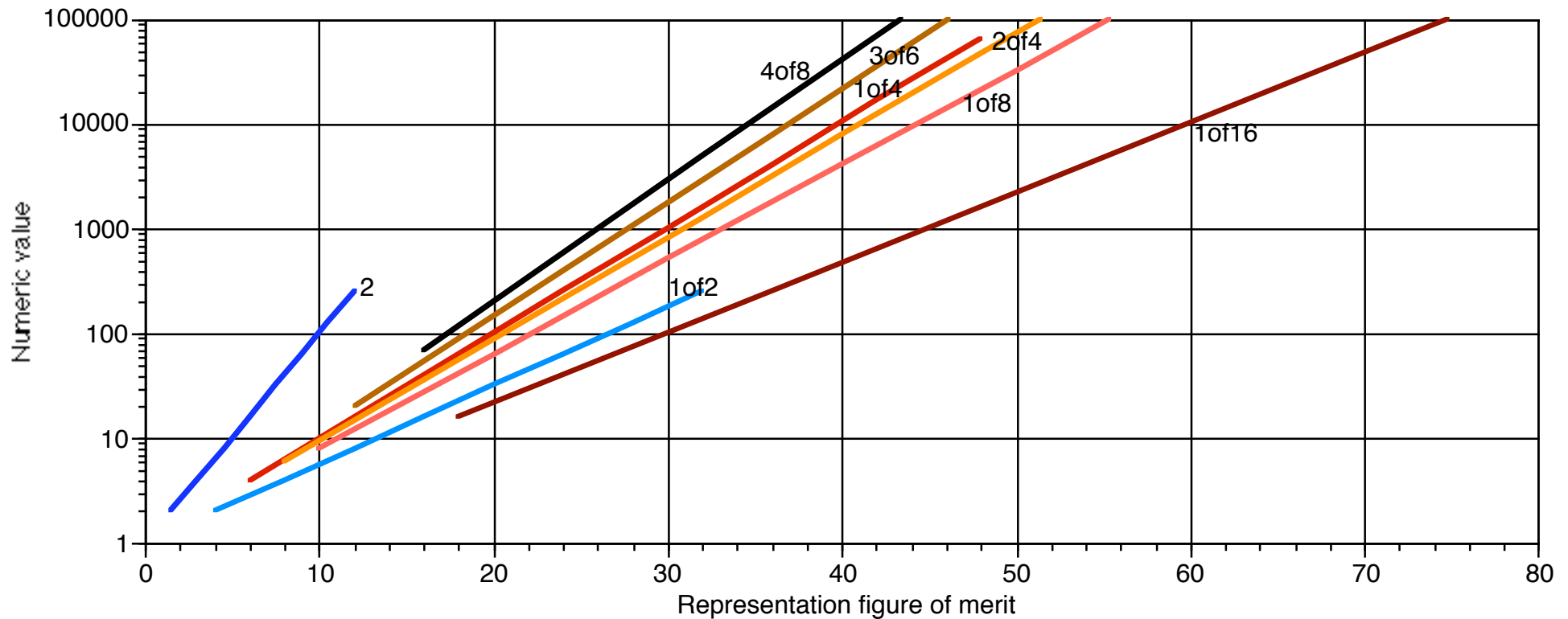
The multi-path encoded digits switch twice per digit because of the NULL transition between each data wavefront. Ignoring the cost of the clock signal and explicit registers, the binary representation switches 0.5 times per digit because of the average switching behavior of binary representation.

The important observation is that 1of4 is better than any of the  $N/2 \text{ of } N$  encodings

# Combined Cost of Representation

Adding the two costs to get a figure of merit. One can see that 1of4 is the optimal 1 of N encoding. It is superior to 1of2. Only 4of8 and 3of6 appear to be better.

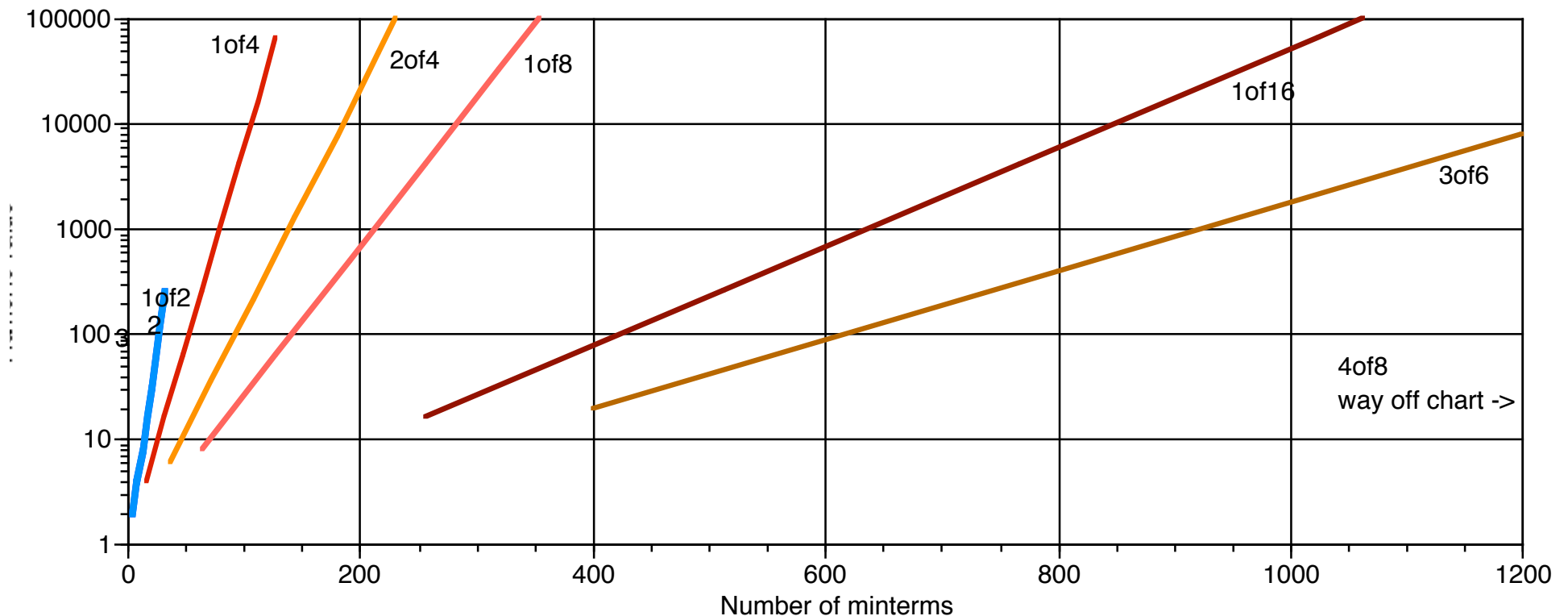
1of4 and 1of2 have the identical representation cost in number of paths but 1of4 has half the switching cost.



# Cost of Expressing Combination

	binary	1of2	1of4	2of4	1of8
The cost of combining digits	4	4	16	36	64

The cost of combining numbers represented in the various encodings with arithmetic operations such as addition is measured by taking the number of minterms defining a function combining two digits.



Again, 1of4 stands out as an optimal representation only bested by 1of2.

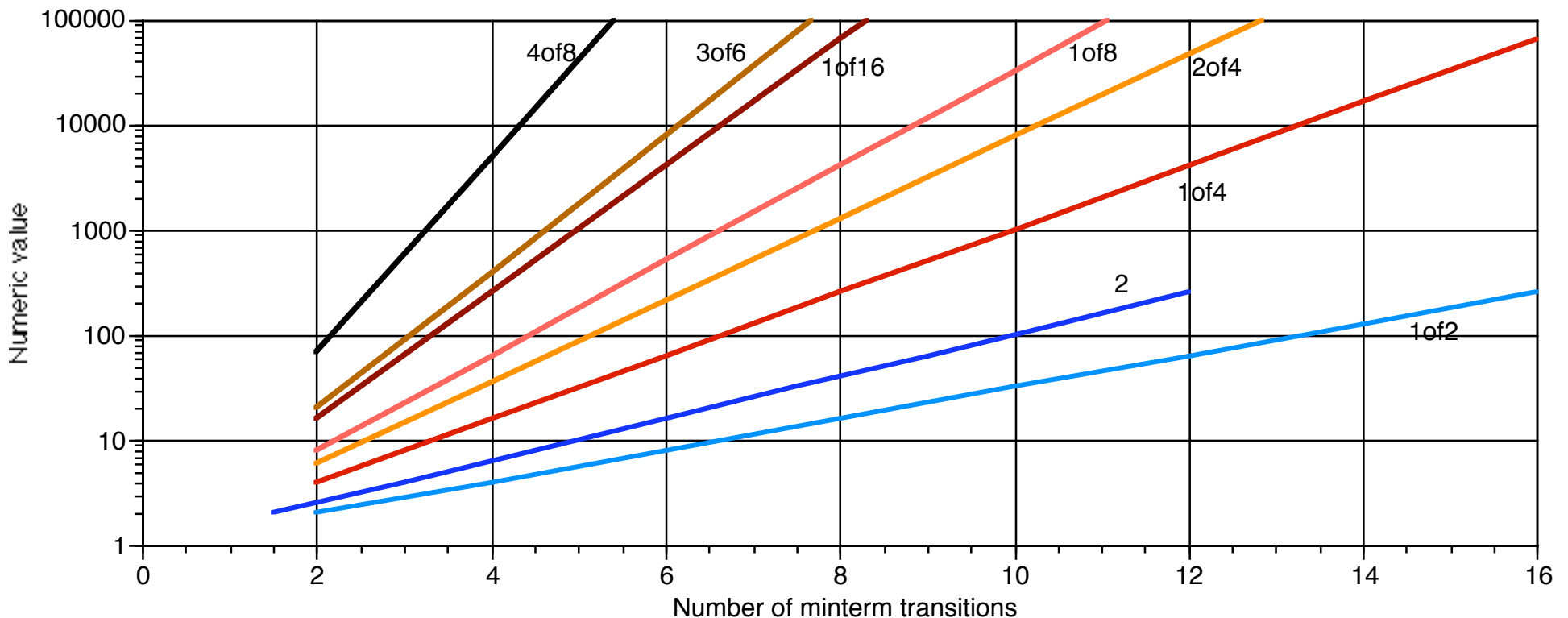
# Cost of Resolving Combination

To represent 64K  
 1 of 2 is  $2 \times 16 = 32$   
 1 of 4 is  $2 \times 8 = 16$   
 1 of 16 is  $2 \times 4 = 8$

The cost of resolving digit combination

	Binary	N/2 of N	1 of N
	1.5	2	2

The cost of resolving digit combination is the number of minterm transitions to resolve a combination of two digits. For all the multi-path encodings exactly one minterm per digit combination will transition to DATA then to NULL for two transitions.

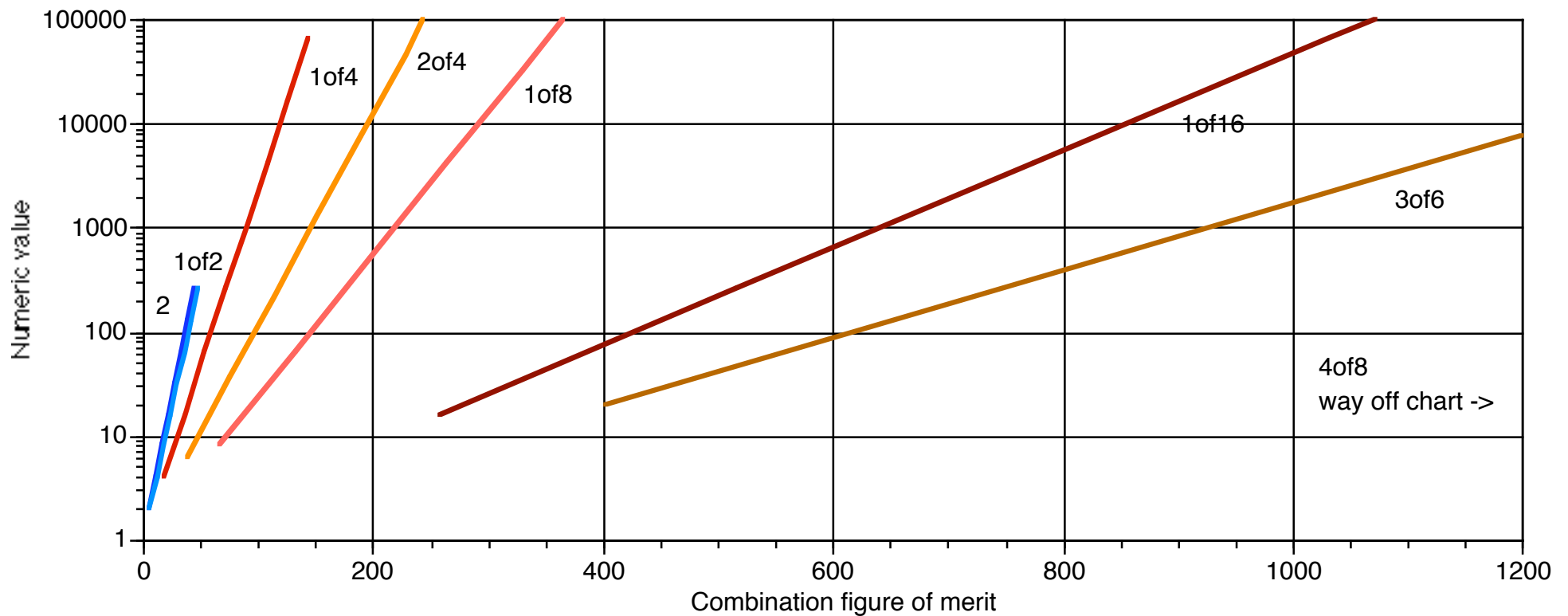


For Boolean representation there are 4 minterms per digit. The probability that the same minterm will be used consecutively and not cause 2 switches is 0.25 leaving an average of 1.5 transitions per digit combination.

The chart does not take into account the fan-in and fan-out complexity of the minterm which is much greater for N/2 of N encodings than for the 1 of N encodings.

# Combined Cost of Combination

When the resource and energy costs of digit combination are combined it shows that 1 of 2 and 1 of 4 are the most efficient encodings.





# Summary of Multi-path Numeric Representation

The 1of4 encoding is clearly superior to 1of2 encoding for transmission. The path resources for 1of2 and 1of4 are identical and 1of4 requires half the switching energy.

The  $N/2$ of $N$  encodings provide slightly better performance for transmission but their combinational performance is abysmal.

The 1of $N > 4$  codes fall below 1of2 and 1of4 in both categories so that leaves only 1of2 and 1of4 to consider.

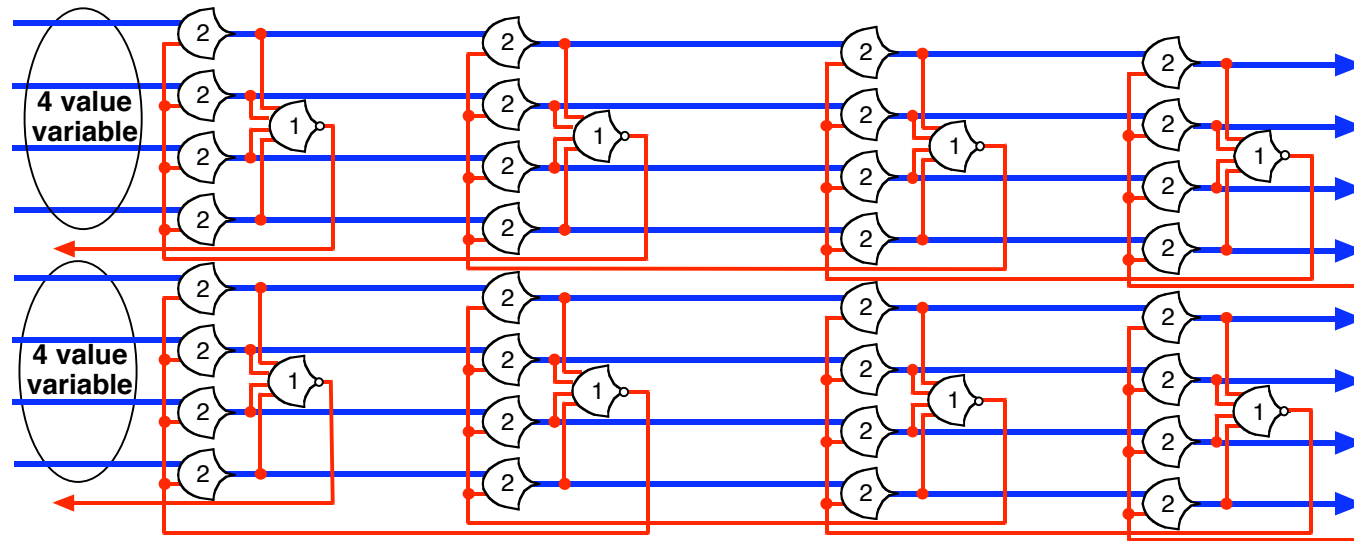
While combining 1of4 variables is more costly in resources, it has a considerable advantage in terms of energy and speed (fewer stages of logic, fewer digits asserting a value, shorter addition carry chain).

The 1of4 (quaternary) encoding seems to be a viable option for numeric representation with multi-path representation.

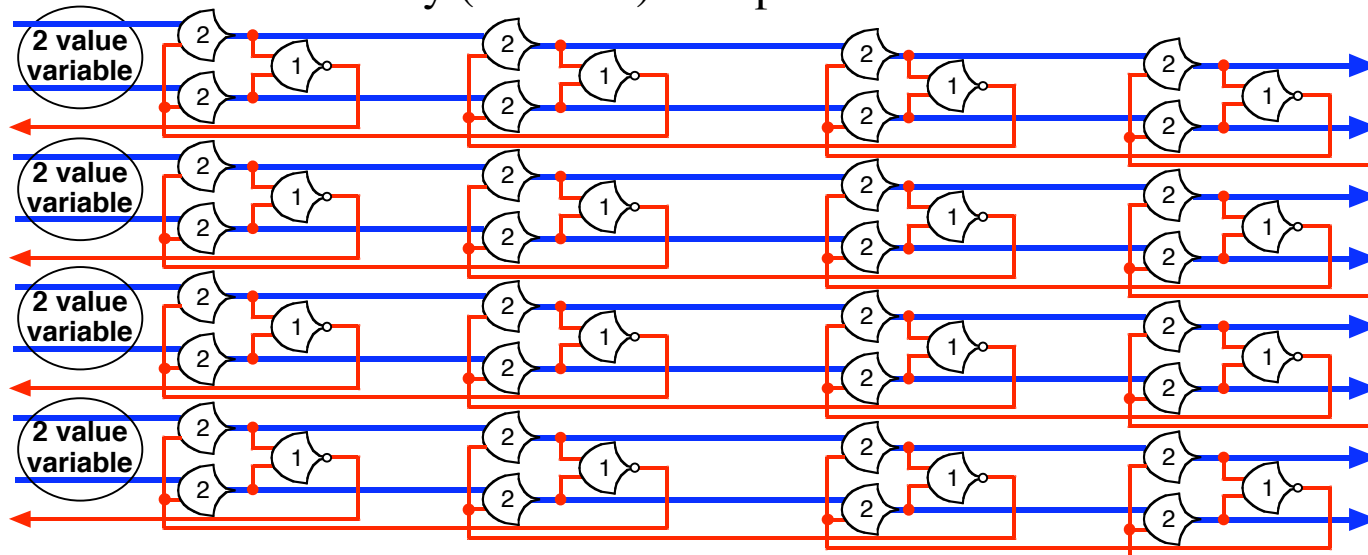
# The Quaternary Data Path

The quaternary data path requires half the switching than does a binary data path with slightly less hardware.

Quaternary (four-rail) data path



Binary (dual-rail) data path



A quaternary ALU and a binary ALU are compared.

The binary ALU is considered both

with 1 of 4 data path representation  
converted to binary and

with 1 of 2 data path representation.

# The Quaternary Logic Operations are Derived in Terms of Tandem Binary Logic Operations.

AND

2 bit binary encoding

	A
	00 01 10 11
00	00 00 00 00
01	00 01 00 01
B 10	00 00 10 10
11	00 01 10 11

Four rail encoding

	A
	0 1 2 3
0	0 0 0 0
B 1	0 1 0 1
2	0 0 2 2
3	0 1 2 3

A0	A1	A2	A3
B0			
B1			
B2			
B3			

OR

2 bit binary encoding

	A
	00 01 10 11
00	00 01 10 11
01	01 01 11 11
B 10	10 11 10 11
11	11 11 11 11

Four rail encoding

	A
	0 1 2 3
0	0 1 2 3
B 1	1 1 3 3
2	2 3 2 3
3	3 3 3 3

A0	A1	A2	A3
B0			
B1			
B2			
B3			

XOR

2 bit binary encoding

	A
	00 01 10 11
00	00 01 10 11
01	01 00 11 10
B 10	10 11 00 01
11	11 10 01 00

Four rail encoding

	A
	0 1 2 3
0	0 1 2 3
B 1	1 0 3 2
2	2 3 0 1
3	3 2 1 0

A0	A1	A2	A3
B0			
B1			
B2			
B3			

NOT

2 bit binary encoding

	A
	00 01 10 11
	11 10 01 00

Four rail encoding

	A
	0 1 2 3
	3 2 1 0

A0	A1	A2	A3

# The Quaternary Arithmetic Operations are Simply Radix Four Arithmetic Operations.

CI = 0

quaternary SUM					
		A			
		0	1	2	3
B	0	0	1	2	3
	1	1	2	3	0
	2	2	3	0	1
	3	3	0	1	2

SUM

	A0	A1	A2	A3
B0				
B1				
B2				
B3				

quaternary CARRY					
		A			
		0	1	2	3
B	0	0	0	0	0
	1	0	0	0	1
	2	0	0	1	1
	3	0	1	1	1

CARRY

	A0	A1	A2	A3
B0				
B1				
B2				
B3				

CI = 1

quaternary SUM					
	A				
	0	1	2	3	
B	0	1	2	3	0
	1	2	3	0	1
	2	3	0	1	2
	3	0	1	2	3

SUM

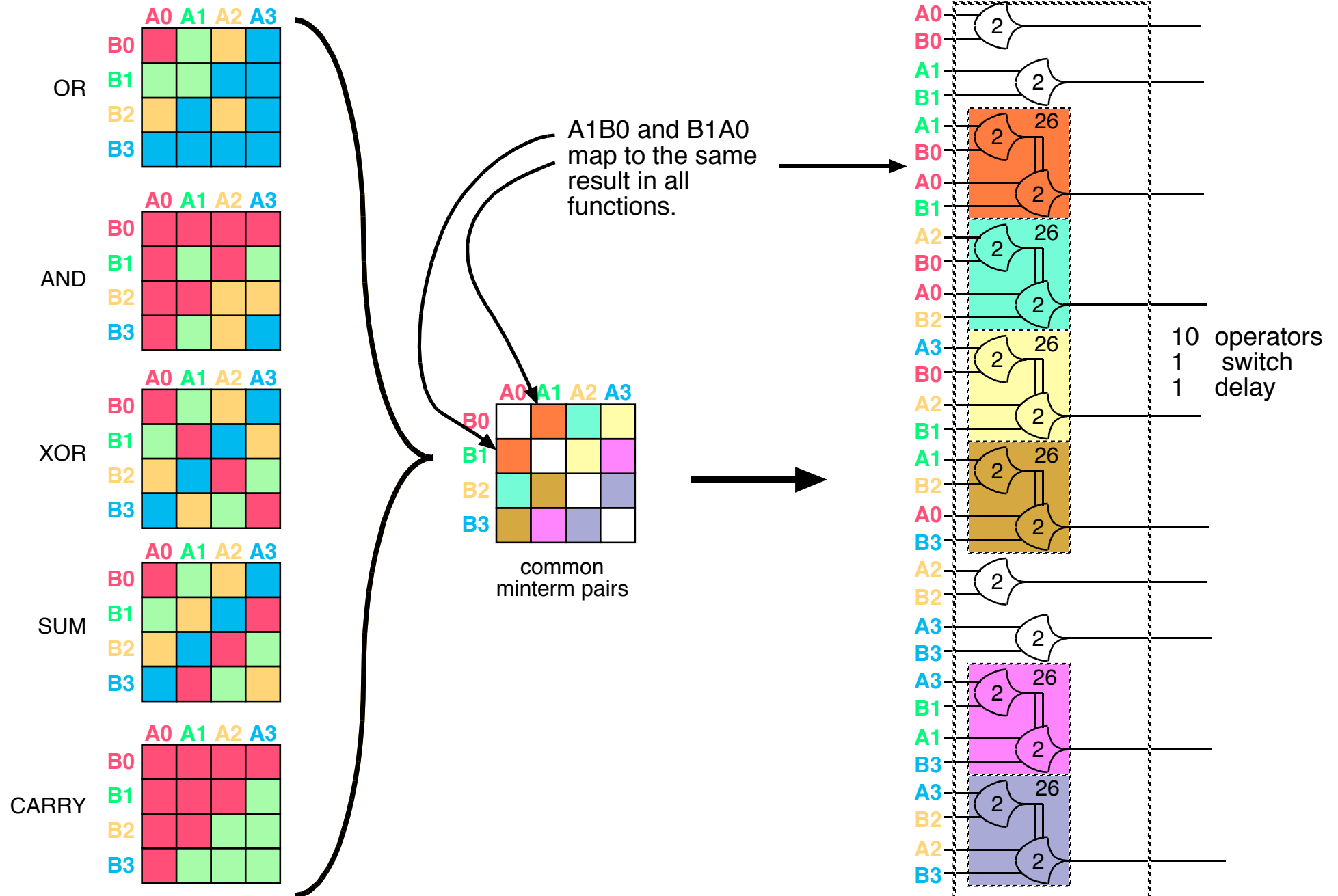
	A0	A1	A2	A3
B0				
B1				
B2				
B3				

quaternary CARRY					
		A			
		0	1	2	3
B	0	0	0	0	1
	1	0	0	1	1
	2	0	1	1	1
	3	1	1	1	1

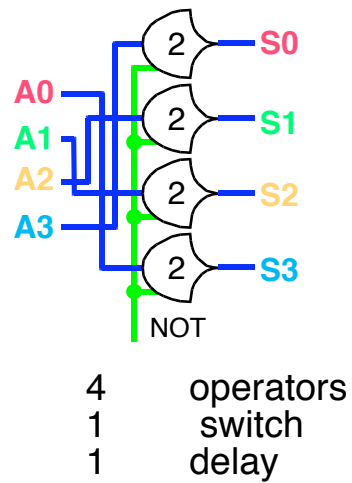
CARRY

	A0	A1	A2	A3
B0				
B1				
B2				
B3				

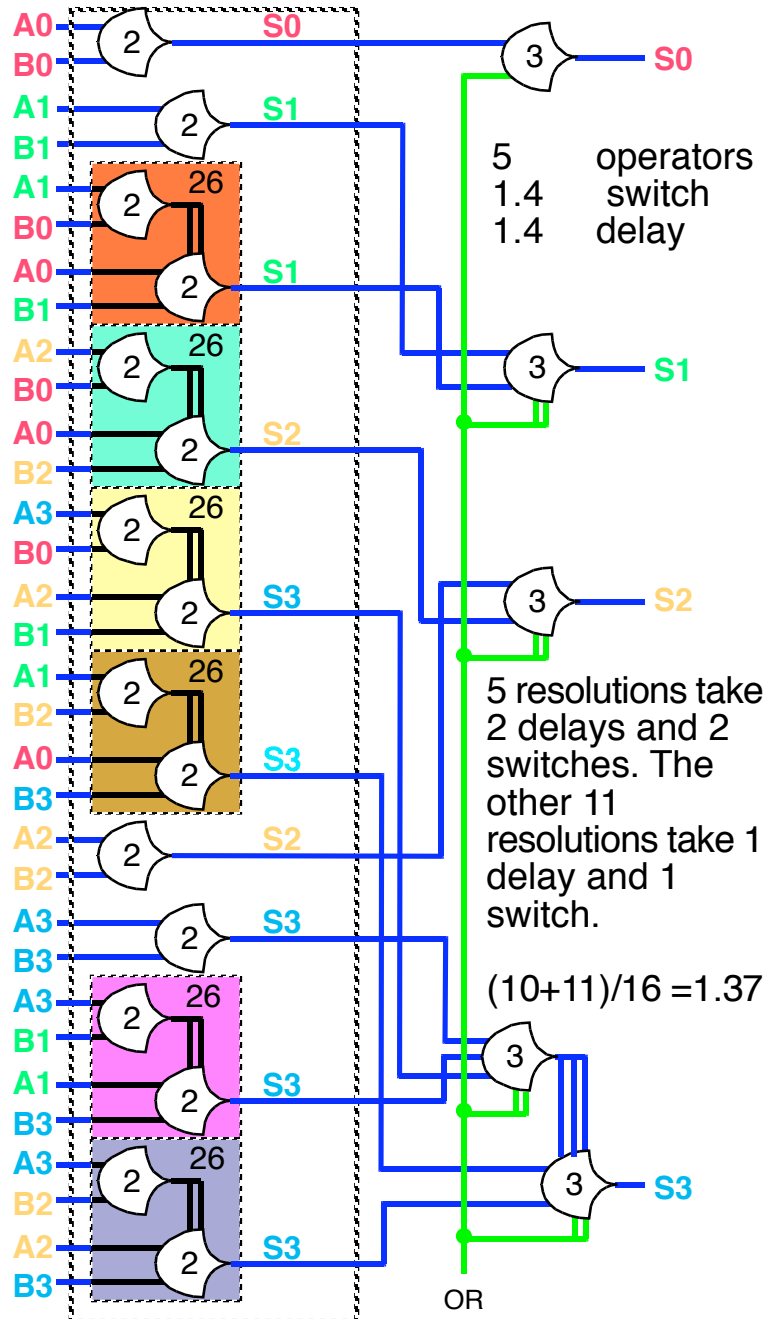
# Common Minterm



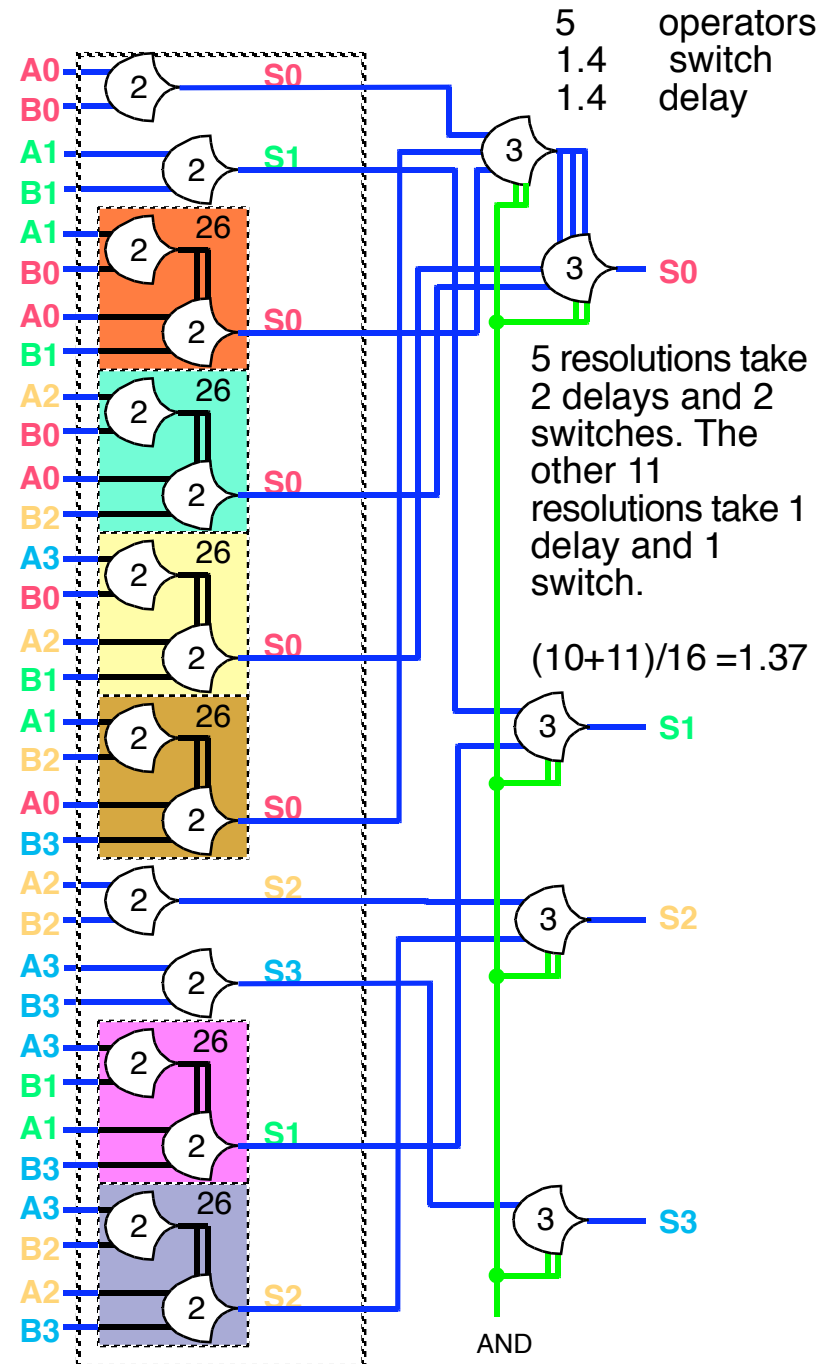
## Quaternary NOT



## Quaternary OR

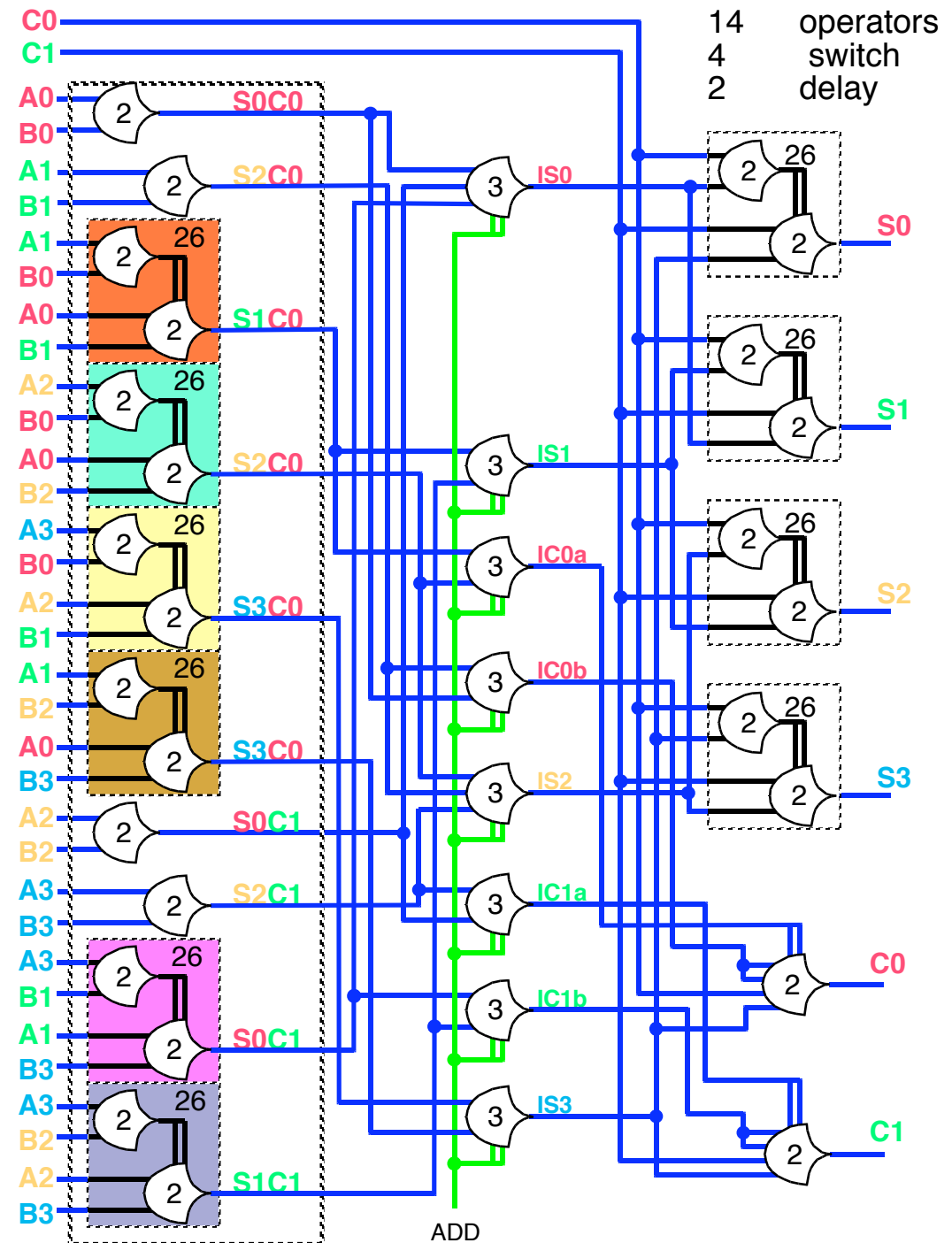
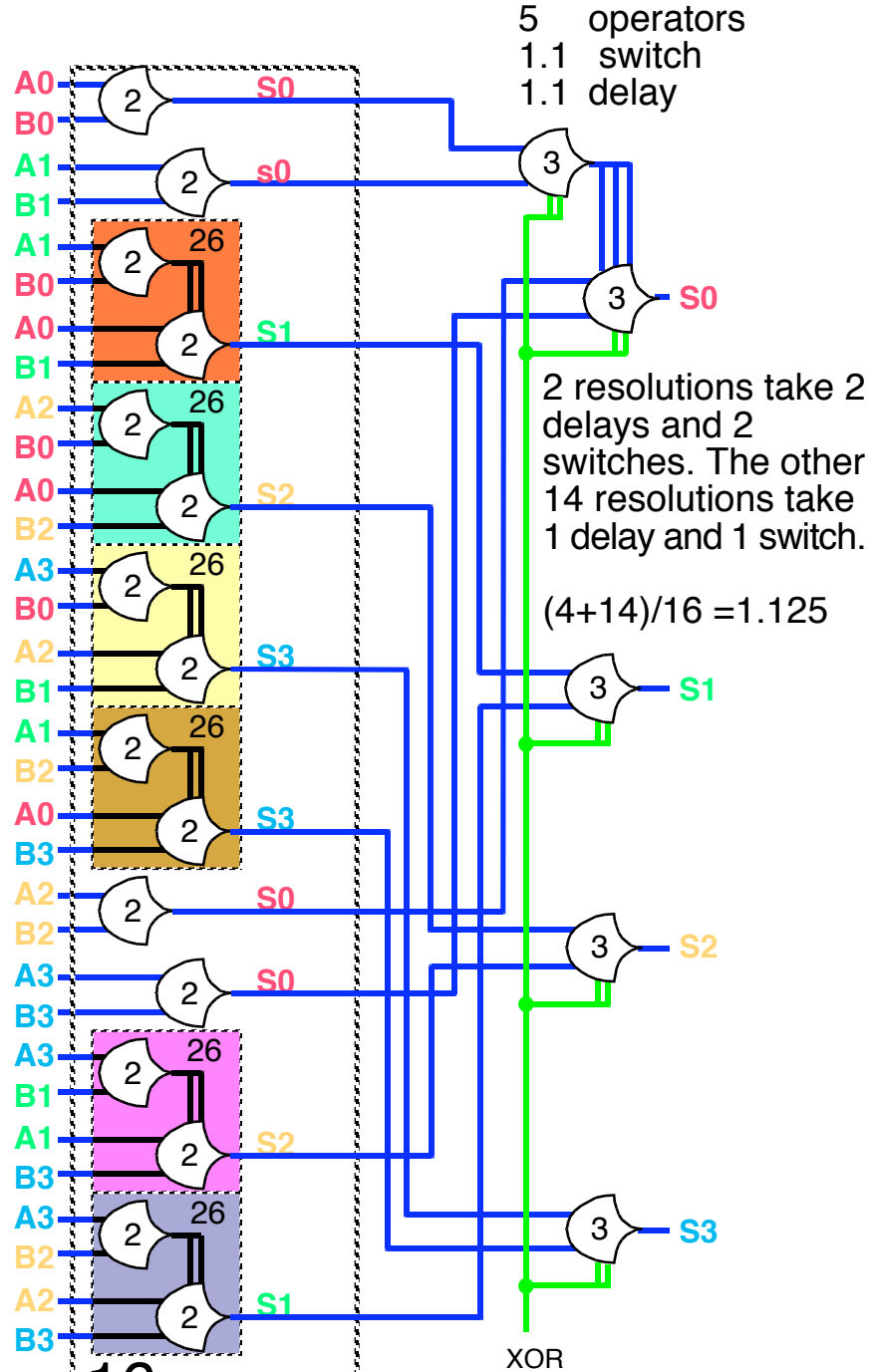


## Quaternary AND



## Quaternary XOR

## Quaternary ADD





# The Quaternary Shift Operation is Derived in Terms of Tandem Binary Shift Operations.

Shift Right  
one bit with  
carry in zero  
(SRC0)

in			out		
carry	quat	binary	binary	quat	carry
0	= 0	= 00	= 00	= 0	= 0
0	= 1	= 01	= 00	= 0	= 1
0	= 2	= 10	= 01	= 1	= 0
0	= 3	= 11	= 01	= 1	= 1



Shift Right  
one bit with  
carry in one  
(SRC1)

in			out		
carry	quat	binary	binary	quat	carry
1	= 0	= 00	= 10	= 2	= 0
1	= 1	= 01	= 10	= 2	= 1
1	= 2	= 10	= 11	= 3	= 0
1	= 3	= 11	= 11	= 3	= 1



Shift Left  
one bit with  
carry in zero  
(SLC0)

in			out		
carry	quat	binary	binary	quat	carry
0	= 0	= 00	= 00	= 0	= 0
0	= 1	= 01	= 10	= 2	= 0
0	= 2	= 10	= 00	= 0	= 1
0	= 3	= 11	= 10	= 2	= 1

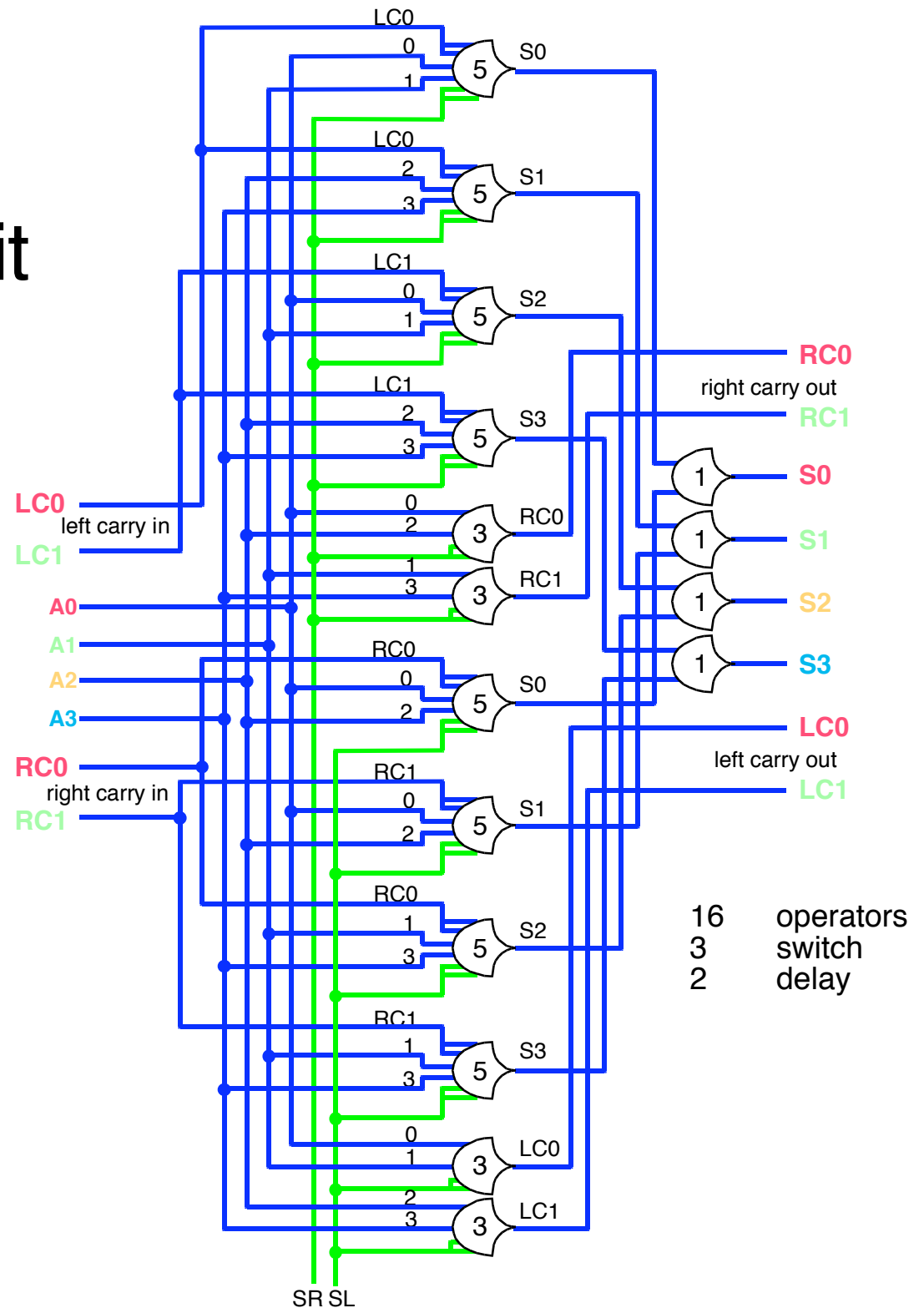


Shift Left  
one bit with  
carry in one  
(SLC1)

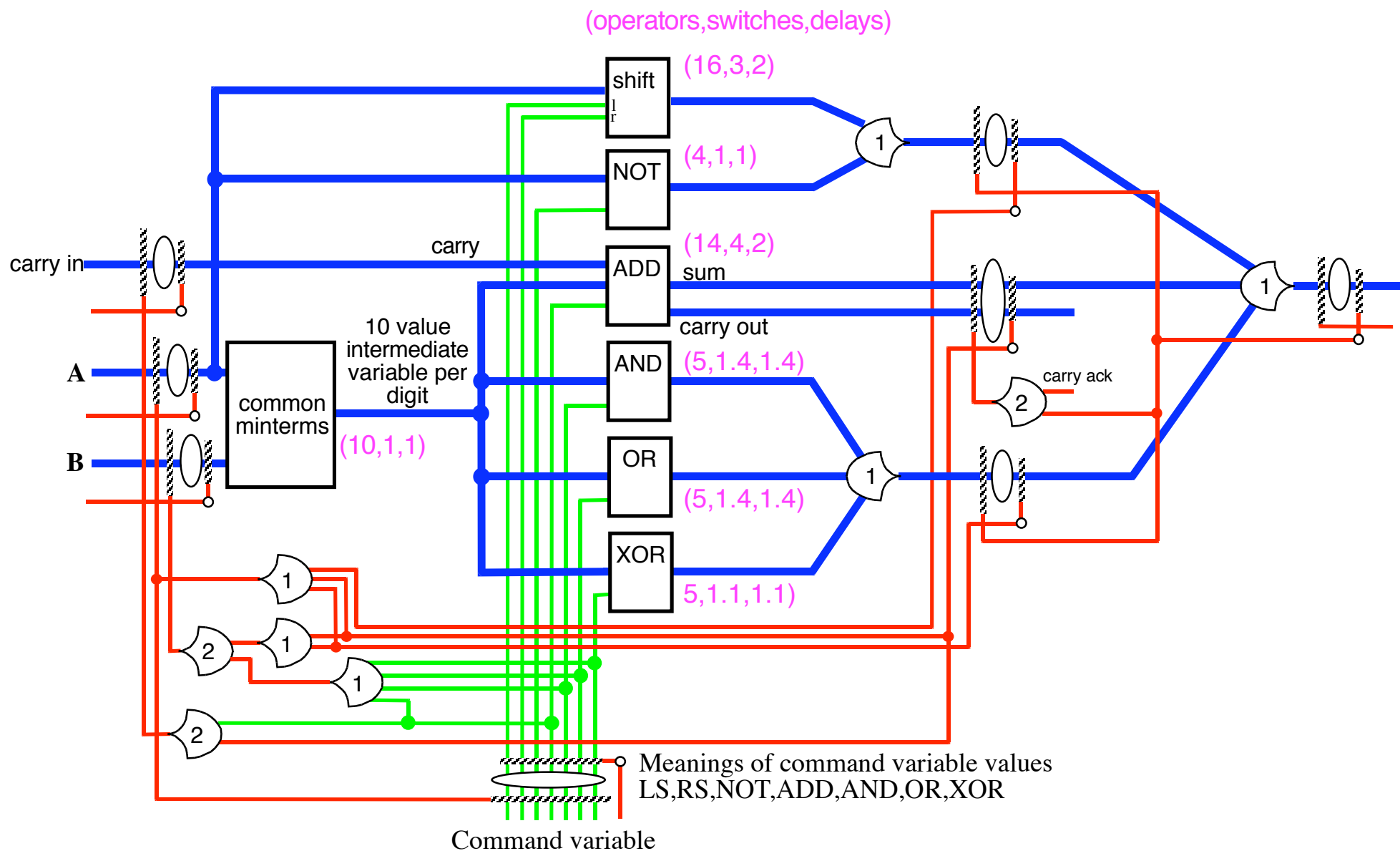
in			out		
carry	quat	binary	binary	quat	carry
1	= 0	= 00	= 01	= 1	= 0
1	= 1	= 01	= 11	= 3	= 0
1	= 2	= 10	= 01	= 1	= 1
1	= 3	= 11	= 11	= 3	= 1



# Quaternary One Bit Shift Left or Right

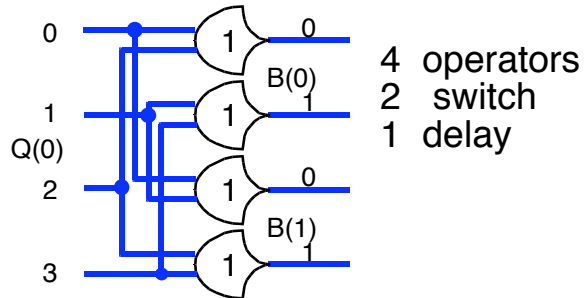


# Quaternary ALU

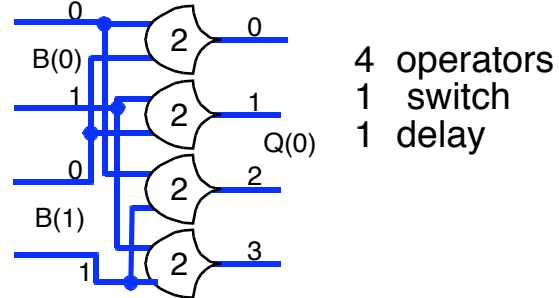


# Binary Operators

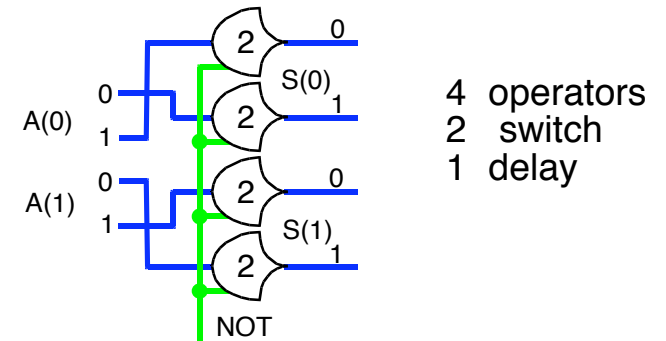
## Quaternary to Binary conversion



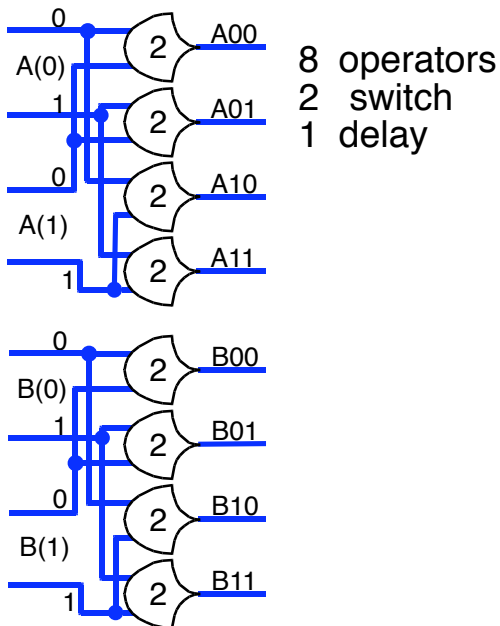
## Binary to Quaternary conversion



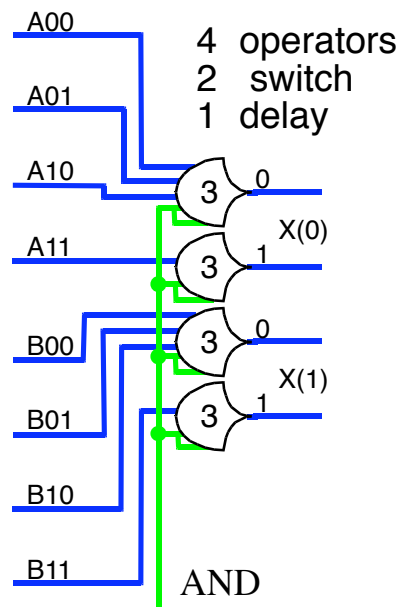
## Binary NOT



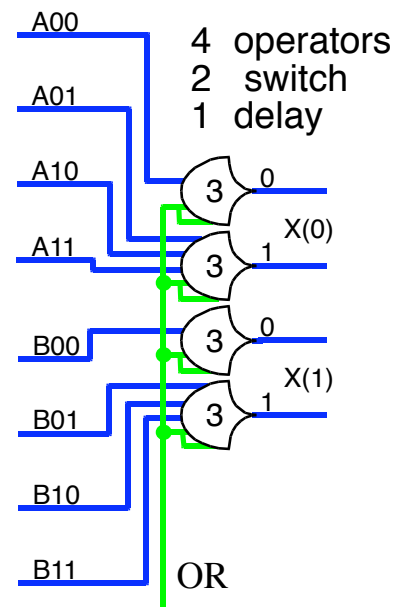
## Common minterm



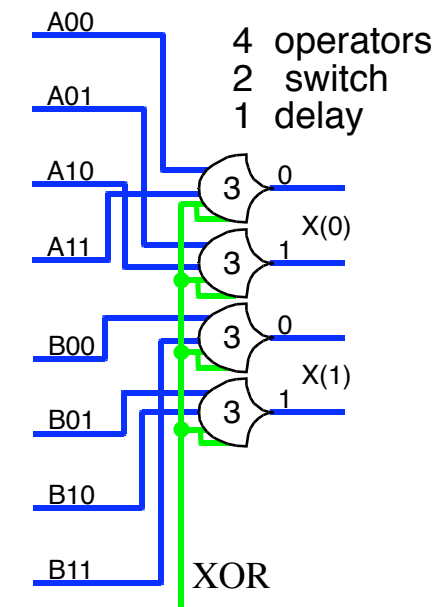
## Binary AND



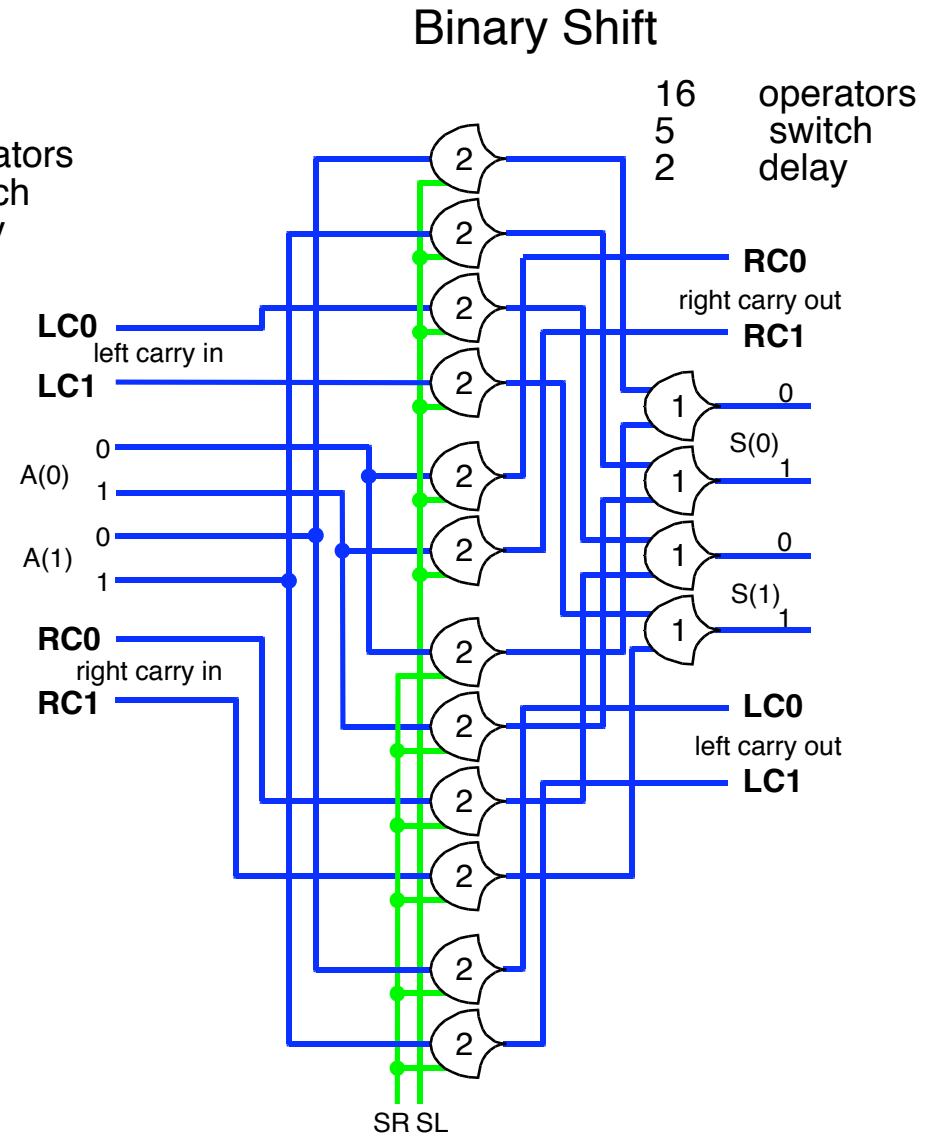
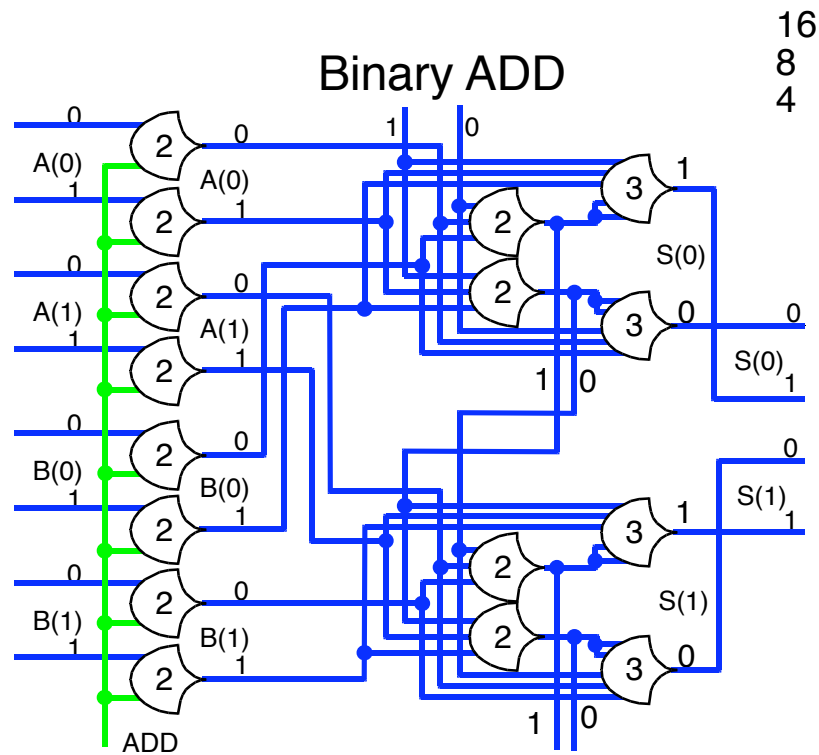
## Binary OR



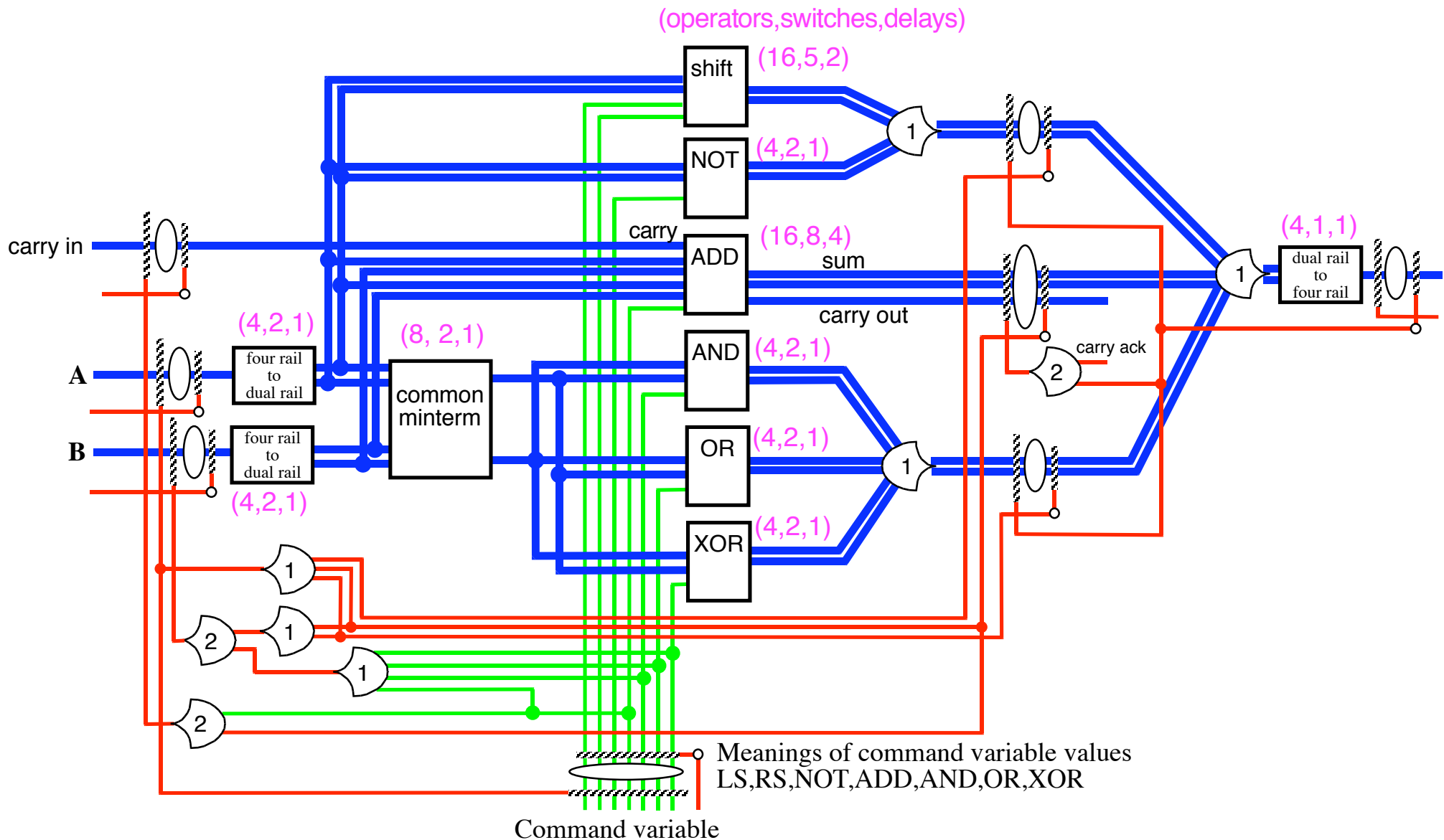
## Binary XOR



# Binary Operators



# Binary ALU With Converison



# ALU Comparisons

The Quaternary ALU is clearly competitive.

Converting from 1 of 4 to 1 of 2 and back is clearly expensive. It is much more economic to remain in 1 of 4 mode and perform quaternary arithmetic.

Is it more economical to remain entirely in 1 of 2 mode?

Removing the data path conversion from the binary ALU removes 2 delays from each operation, 5 switchings from AND, OR XOR and ADD and 3 switchings from NOT and SHIFT. The resulting ALU has slightly fewer operators but is still more expensive in terms of switching and delay and the switching cost of data path transmission doubles.

	<i>Binary ALU with conv</i>			<i>Quaternary ALU</i>			<i>Binary ALU without conv</i>		
	<i>operators</i>	<i>switches</i>	<i>delay</i>	<i>operators</i>	<i>switches</i>	<i>delay</i>	<i>operators</i>	<i>switches</i>	<i>delay</i>
4 to 2	8	4	1						
2 to 4	4	1	1						
com min	8	1	1	10	1	1	8	1	1
ADD	16	13	6	14	5	3	16	8	4
OR	4	9	4	5	2.4	2.4	4	4	2
AND	4	9	4	5	2.4	2.4	4	4	2
XOR	4	9	4	5	2.1	2.1	4	4	2
NOT	4	5	3	4	1	1	4	2	1
SHIFT	16	8	4	16	3	2	16	5	2
<b>totals</b>	<b>68</b>	<b>53</b>	<b>25</b>	<b>59</b>	<b>15.9</b>	<b>12.9</b>	<b>56</b>	<b>27</b>	<b>13</b>

The red entries indicate differences with the table in the book. Apparently I did not vet the table as carefully as I thought I had.

The differences are minor and the conclusion remains the same.

For comparison purposes the raw totals are shown instead of dividing all the switch and delay totals by 6 to get an average per operator.

The switch and delay numbers for the operations include the conversions and minterm as a total cost of each operator.

This page intentionally blank